## Bedford-McMullen carpets - an example in dimension theory of self-affine sets

In this note we study the dimension of the classical Bedford-McMullen carpets. We apply the learned methods from the mini-course of Károly Simon on "Dimension Theory of self-affine and almost self-affine sets and measures" during Simons Semester in Banach Center, October 2015.

Let $n>m \geq 2$ be integers and let $A \subseteq\{0, \ldots, m-1\} \times\{0, \ldots, n-1\}$. We consider the following subset of $[0,1]^{2}$

$$
\Lambda:=\left\{\binom{\sum_{k=1}^{\infty} \frac{i_{k}}{m^{k}}}{\sum_{k=1}^{\infty} \frac{j_{k}}{n^{k}}}:\left(i_{k}, j_{k}\right) \in A \text { for every } k \geq 1\right\}
$$

It is easy to see that $\Lambda$ is a self-affine set on the plane, i.e.

$$
\Lambda=\bigcup_{(i, j) \in A} f_{(i, j)}(\Lambda), \text { where } f_{(i, j)}\binom{x}{y}=\left(\begin{array}{cc}
\frac{1}{m} & 0 \\
0 & \frac{1}{n}
\end{array}\right)\binom{x}{y}+\binom{\frac{i}{m}}{\frac{j}{n}}
$$



Introduce some notations:

$$
\begin{aligned}
& M:=\sharp A \\
& Q:=\{i: \exists j \text { s.t. }(i, j) \in A\} \\
& N:=\sharp Q \\
& T_{i}:=\{j:(i, j) \in A\} \\
& t_{i}:=\sharp T_{i}
\end{aligned}
$$

First, let us observe that the projection of $\Lambda$ to the vertical axes is a self-similar set on the line. Precisely, $\operatorname{proj} \Lambda$ is the attractor of the IFS

$$
\begin{equation*}
\left\{g_{i}: x \mapsto \frac{x+i}{m}\right\}_{i \in Q} \tag{1}
\end{equation*}
$$

where proj denotes the orthogonal projection to the $x$-axes.

Falconer showed that the Hausdorff and box dimension of any self-similar set are equal, moreover, since $\left\{g_{i}: x \mapsto \frac{x+i}{m}\right\}_{i \in Q}$ satisfies the open set condition w.r.t the interval $(0,1)$ we get

$$
\begin{equation*}
\operatorname{dim}_{H} \operatorname{proj} \Lambda=\operatorname{dim}_{B} \operatorname{proj} \Lambda=\frac{\log N}{\log m} \tag{2}
\end{equation*}
$$

Our first approach is to calculate the root of the subadditive pressure, i.e. the affinity dimension. Let us recall the definition of the pressure function. Let

$$
\phi^{s}(A):= \begin{cases}\alpha_{1}(A)^{s} & \text { if } 0 \leq s<1 \\ \alpha_{1}(A) \alpha_{2}(A)^{s-1} & \text { if } 1 \leq s<2 \\ \left(\alpha_{1}(A) \alpha_{2}(A)\right)^{s / 2} & \text { if } s \geq 2\end{cases}
$$

where $\alpha_{i}(A)$ denotes the $i$ th singular value of the matrix $A$. Then let $s_{0}$ be the unique root of the strictly monotone decreasing and continuous function $s \mapsto P(s)$, where

$$
P(s)=\lim _{n \rightarrow \infty} \frac{1}{n} \log \sum_{\left(i_{1}, j_{1}\right), \ldots,\left(i_{n}, j_{n}\right) \in A} \phi^{s}\left(f_{\left(i_{1}, j_{1}\right)}^{\prime} \cdots f_{\left(i_{n}, j_{n}\right)}^{\prime}\right)
$$

the subadditive pressure function. Because of the special form of our matrices, one can show that

$$
\begin{equation*}
s_{0}=\min \left\{\frac{\log M}{\log m}, 1+\frac{\log M-\log m}{\log n}\right\} \tag{3}
\end{equation*}
$$

Falconer showed that this is an upper bound for the box counting dimension for every self-affine set.

Box counting dimension. First, let us define the lower and upper box counting dimension of a bounded set $A \subset \mathbb{R}^{2}$.

$$
\underline{\operatorname{dim}}_{B} A=\liminf _{\delta \rightarrow 0+} \frac{\log N_{\delta}(A)}{-\log \delta} \text { and } \overline{\operatorname{dim}}_{B} A=\limsup _{\delta \rightarrow 0+} \frac{\log N_{\delta}(A)}{-\log \delta}
$$

where $N_{\delta}(A)=\min \left\{N: \exists \underline{x}_{1}, \ldots, \underline{x}_{N} \in \mathbb{R}^{2}\right.$ s.t. $\left.A \subseteq \bigcup_{i=1}^{N} B_{\delta}\left(\underline{x}_{i}\right)\right\}$ and $B_{\delta}(\underline{x})$ denotes the closed ball with radius $\delta$ and centered at $\underline{x}$. If the limit exists then we denote it by $\operatorname{dim}_{B} A$.

Let us observe that in the liminf and limsup $\delta$ can be replaced by any exponential sequence. That is,

$$
\underline{\operatorname{dim}}_{B} A=\liminf _{l \rightarrow \infty} \frac{\log N_{1 / n^{l}}(A)}{l \log n} \text { and } \overline{\operatorname{dim}}_{B} A=\limsup _{l \rightarrow \infty} \frac{\log N_{1 / n^{l}}(A)}{l \log n} .
$$

Now, we construct an optimal cover for $\Lambda$. Note that for every $l \geq 1$ and every $\left(i_{1}, j_{1}\right), \ldots,\left(i_{l}, j_{l}\right)$ the rectangles $f_{\left(i_{1}, j_{1}\right)} \circ \cdots \circ f_{\left(i_{l}, j_{l}\right)}\left((0,1)^{2}\right)$ are disjoint. Let $k \geq 1$ the smallest integer such that $1 / m^{k} \leq 1 / n^{l}$, i.e. $k=\left\lceil l \frac{\log n}{\log m}\right\rceil$. Divide the horizontal side (which has length $m^{-l}$ ) of the rectangle $f_{\left(i_{1}, j_{1}\right)} \circ \cdots \circ f_{\left(i_{l}, j_{l}\right)}\left((0,1)^{2}\right)$ into $m^{k-l}$ equal parts. Denote these "approximate squares" from left to right by $R_{1}^{\mathbf{i}}, R_{2}^{\mathbf{i}}, \ldots, R_{m^{k-l}}^{\mathbf{i}}$, where $\mathbf{i}=\left(i_{1}, j_{1}\right), \ldots,\left(i_{l}, j_{l}\right)$. Thus, any of the rectangles has vertical side length $n^{-l}$ and horizontal side length $m^{-k}$.


In the general case, for the definition of subadditive pressure we used all of the approximate squares to cover the set. In the case of Bedford-McMullen carpet, because of the special structure, we do not need all of them. By applying the inverse function $f_{\left(i_{l}, j_{l}\right)}^{-1} \circ \cdots \circ f_{\left(i_{1}, j_{1}\right)}^{-1}$ for the rectangle $f_{\left(i_{1}, j_{1}\right)} \circ \cdots \circ f_{\left(i_{l}, j_{l}\right)}\left((0,1)^{2}\right)$ we get the following picture:

where $\widetilde{R}_{k}^{\mathbf{i}}=f_{\left(i_{l}, j_{l}\right)}^{-1} \circ \cdots \circ f_{\left(i_{1}, j_{1}\right)}^{-1}\left(R_{k}^{\mathbf{i}}\right)$. So we see that we may choose only those $\widetilde{R}_{k}^{\mathbf{i}}$ columns for which $\widetilde{R}_{k}^{\mathbf{i}} \cap \Lambda \neq \emptyset$. The number of such $\widetilde{R}_{k}^{\mathbf{i}}$ columns is equal to the number of non-empty columns in the $k-l$ th iteration. Or in other words, the number of intervals with lenght $m^{l-k}$ needed to cover proj $\Lambda$. So by (2), for every $\mathbf{i} \in A^{l}$

$$
\sharp\left\{R_{k}^{\mathrm{i}}: R_{k}^{\mathrm{i}} \cap \Lambda \neq \emptyset \text { for } k=1, \ldots, m^{k-l}\right\}=N^{k-l} .
$$

Therefore

$$
\widetilde{N}_{l}:=\left\{R_{k}^{\mathbf{i}}: \mathbf{i} \in A^{l} \& R_{k}^{\mathbf{i}} \cap \Lambda \neq \emptyset \text { for } k=1, \ldots, m^{k-l}\right\} \text { and } \sharp \widetilde{N}_{l}=M^{l} N^{k-l} \text {. }
$$

Since every rectangle $R_{k}^{\mathbf{i}} \in \widetilde{N}_{l}$ can be extended to a ball with radius $1 / n^{l}$. Thus, $\sharp \widetilde{N}_{l} \geq N_{1 / n^{l}}(\Lambda)$ and

$$
\begin{aligned}
\overline{\operatorname{dim}}_{B} \Lambda \leq \limsup _{l \rightarrow \infty} \frac{\log M^{l} N^{k-l}}{l \log n} & = \\
& \quad \limsup _{l \rightarrow \infty} \frac{\log M}{\log n}+\left(\frac{\lceil l \log n}{\log m}\right\rceil \\
l & 1) \frac{\log N}{\log n}=\frac{\log M}{\log n}+\left(1-\frac{\log m}{\log n}\right) \frac{\log N}{\log m} .
\end{aligned}
$$

On the other hand, let $\mathcal{B}_{l}:=\left\{B_{1 / n^{l}}\left(\underline{x}_{1}\right), \ldots, B_{1 / n^{l}}\left(\underline{x}_{N_{1 / n} l}(\Lambda)\right)\right\}$ be the set of balls which covers optimally the set $\Lambda$. Then for every $R_{k}^{\mathrm{i}} \in \widetilde{N}_{l}$ intersects at least one $B \in \mathcal{B}_{l}$, moreover, any ball $B \in \mathcal{B}_{l}$ may intersect at most $3 m$ approximate squares from $\widetilde{N}_{l}$. Hence,

$$
\sharp \tilde{N}_{l} \leq 3 m N_{1 / n^{l}}(\Lambda)
$$

and therefore,

$$
\underline{\operatorname{dim}}_{B} \Lambda \geq \liminf _{l \rightarrow \infty} \frac{\log (3 m)^{-1} M^{l} N^{k-l}}{l \log n}=\frac{\log M}{\log n}+\left(1-\frac{\log m}{\log n}\right) \frac{\log N}{\log m}
$$

In summary, we get that the box counting dimension exists and

$$
\begin{equation*}
\operatorname{dim}_{B} \Lambda=\frac{\log M}{\log n}+\left(1-\frac{\log m}{\log n}\right) \frac{\log N}{\log m} \tag{4}
\end{equation*}
$$

Remark 1. By using the formulas (3) and (4), simple algebraic calculations show that

$$
s_{0}=\operatorname{dim}_{B} \Lambda \Leftrightarrow N=M \text { or } N=m .
$$

In other words, by (2), the box dimension is equal to the affinity dimension if and only if

$$
\operatorname{dim}_{B} \operatorname{proj} \Lambda=\min \left\{1, \operatorname{dim}_{B} \Lambda\right\} .
$$

Lower bound for Hausdorff dimension. Now, we turn to the case of Hausdorff dimension. Let $\mu$ be a self-affine measure with probability vector $\underline{p}=\left(p_{(i, j)}\right)_{(i, j) \in A}$. That is, $\mu$ is the unique compactly supported measure, for which

$$
\int h(x) d \mu(x)=\sum_{(i, j) \in A} p_{(i, j)} \int h\left(f_{(i, j)}(x)\right) d \mu(x),
$$

for any continuous test function $h$ on $\Lambda$.
Let us observe again, that the projection of the measure $\mu$ onto the $x$-axes is a self-similar measure (like in the case of the set). That is,

$$
\operatorname{proj}_{*} \mu=\sum_{i \in Q}\left(\sum_{j \in T_{i}} p_{i, j}\right) \operatorname{proj}_{*} \mu \circ g_{i}^{-1},
$$

where $\operatorname{proj}_{*} \mu=\mu \circ \operatorname{proj}^{-1}$ and $g_{i}$ s are from the IFS (1).
By the Feng-Hu formula, we are able to calculate the Hausdorff dimension of the measure $\mu$, i.e.

$$
\begin{equation*}
\operatorname{dim}_{H} \mu=\frac{-\sum_{(i, j) \in A} p_{i, j} \log p_{i, j}}{\log n}+\left(1-\frac{\log m}{\log n}\right) \frac{-\sum_{i \in Q} \sum_{j \in T_{i}} p_{i, j} \log \sum_{j \in T_{i}} p_{i, j}}{\log m} . \tag{5}
\end{equation*}
$$

For simplicity, let us denote $\sum_{j \in T_{i}} p_{i, j}$ by $q_{i}$.
Remark 2. The formulas (4) and (5) are very similar to each other. By the definition of the entropy,

$$
\log M=-\sum_{(i, j) \in A} p_{i, j} \log p_{i, j} \Leftrightarrow p_{i, j}=\frac{1}{M}
$$

and

$$
\log N=-\sum_{i \in Q} q_{i} \log q_{i} \Leftrightarrow q_{i}=\frac{1}{N} .
$$

Thus, $\operatorname{dim}_{B} \Lambda=\operatorname{dim}_{H} \mu$ for a probability vector $\underline{p}=\left(p_{(i, j)}\right)_{(i, j) \in A}$ if and only if $p_{i, j}=1 / M$ for every $(i, j) \in A$ and $t_{i}=M / N$ for every $i \in Q$.

By definition, $\operatorname{dim}_{H} \Lambda \geq \operatorname{dim}_{H} \mu$ therefore to get a lower bound, we maximize the value of (5). Use the method of Lagrange-multipliers! That is, we maximize the function

$$
d(\underline{p}, \lambda)=\operatorname{dim}_{H} \mu+\lambda\left(\sum_{(i, j) \in A} p_{i, j}-1\right) .
$$

It is easy to see that $d(\underline{p}, \lambda)$ is concave. By taking the derivative w.r.t $p_{i, j}$ we get

$$
\frac{-\log p_{i, j}-1}{\log n}+\left(1-\frac{\log m}{\log n}\right) \frac{-\log \sum_{j \in T_{i}} p_{i, j}-1}{\log m}+\lambda=0 .
$$

Thus, $p_{i, j}=q_{i} / t_{i}$ for every $(i, j) \in A$ (for fixed $i$ it the value is independent of $j$.) Thus, it is enough to maximize the function

$$
(\underline{q}, \lambda) \mapsto \frac{-\sum_{i \in Q} q_{i} \log q_{i} / t_{i}}{\log n}+\left(1-\frac{\log m}{\log n}\right) \frac{-\sum_{i \in Q} q_{i} \log q_{i}}{\log m}+\lambda\left(\sum_{i \in Q} q_{i}-1\right) .
$$

By taking the derivative w.r.t $q_{i}$ we get

$$
\frac{\log q_{i} / t_{i}-1}{\log n}+\left(1-\frac{\log m}{\log n}\right) \frac{-\log q_{i}-1}{\log m}+\lambda=0
$$

Thus,

$$
\begin{equation*}
q_{i}=\frac{t_{i}^{\frac{\log m}{\log n}}}{\sum_{i^{\prime} \in Q} t_{i^{\prime}}^{\frac{\log m}{\log }}} \text { and } p_{i, j}=\frac{t_{i}^{\frac{\log m}{\log n}-1}}{\sum_{i^{\prime} \in Q} t_{i^{\prime}}^{\frac{\log m}{\log n}}} . \tag{6}
\end{equation*}
$$

Hence, we get

$$
\begin{equation*}
\operatorname{dim}_{H} \Lambda \geq \frac{\sum_{i \in Q} t_{i}^{\frac{\log m}{\log n}}}{\log m} \tag{7}
\end{equation*}
$$

Upper bound for Hausdorff dimension. Our claim is that the lower bound in (7) is sharp. One way to show that is to find an optimal cover for the set $\Lambda$. However, our natural cover, which was constructed to calculate the box dimension, is not optimal if there is an $i \in Q$ such that $t_{i} \neq M / N$, see Remark 2. Therefore, we use here a mass distribution principle.
Lemma 1. Let $\nu$ be a probability measure on a set $B \subset \mathbb{R}^{d}$ such that $\nu(B)=1$ and

$$
\liminf _{r \rightarrow 0+} \frac{\log \nu\left(B_{r}(x)\right)}{\log r} \leq \alpha \text { for every } x \in B
$$

where $B_{r}(x)$ denotes the ball centered at $x$ with radius $r$. Then $\operatorname{dim}_{H} B \leq \alpha$.
Proof. Let us recall here the definition of Hausdorff measure, i.e.

$$
\mathcal{H}_{\delta}^{s}(B)=\inf \left\{\sum_{i}\left|U_{i}\right|^{s}: B \subseteq \bigcup_{i} U_{i} \&\left|U_{i}\right|<\delta\right\} \text { and } \mathcal{H}^{s}(B)=\sup _{\delta>0} \mathcal{H}_{\delta}^{s}(B)
$$

By our assumption, for every $\varepsilon, \delta>0$ and every $x \in B$ there exists $\delta>R(x)>0$ such that

$$
\nu\left(B_{R(x)}(x)\right) \geq R(x)^{\alpha+\varepsilon}
$$

Since $\bigcup_{x \in B} B_{R(x)}(x)$ is a cover of $B$, by Besicovitch's covering theorem we get that there exists a $c>0$ and countable subsets $\mathcal{B}_{j}, j=1, \ldots, c$, of the family of balls $\left\{B_{R(x)}(x)\right\}_{x \in B}$ such that

$$
\bigcup_{j=1}^{c} \bigcup_{U \in \mathcal{B}_{j}} U \supseteq B \text { and } U \cap U^{\prime}=\emptyset \text { for every } U \neq U^{\prime} \in \mathcal{B}_{j}
$$

Thus,

$$
\mathcal{H}_{\delta}^{\alpha+\varepsilon}(B) \leq \sum_{j=1}^{c} \sum_{U \in \mathcal{B}_{j}}|U|^{\alpha+\varepsilon} \leq \sum_{j=1}^{c} \sum_{U \in \mathcal{B}_{j}} \nu(U)=c \text { and therefore } \mathcal{H}^{\alpha+\varepsilon}(B) \leq c
$$

Since $\varepsilon>0$ was arbitrary, the statement follows.
We apply Lemma 1 for the self-affine measure $\mu$ with probability vector defined in (6). Let $x=\left(\sum_{k=1}^{\infty} \frac{i_{k}}{m^{k}}, \sum_{k=1}^{\infty} \frac{j_{k}}{n^{k}}\right)^{T} \in \Lambda$ and let $l \geq 1$ integer. Denote by $C_{l}(x)$ the following approximate square

$$
C_{l}(x)=\left\{\binom{\sum_{r=1}^{\infty} \frac{i_{r}^{\prime}}{m^{r}}}{\sum_{r=1}^{\infty} \frac{j_{r}^{\prime}}{n^{r}}} \in \Lambda: i_{p}=i_{p}^{\prime} \text { for } p=1, \ldots, k \text { and } j_{q}=j_{q}^{\prime} \text { for } q=1, \ldots, l\right\}
$$

where $k=\left\lceil l \frac{\log n}{\log m}\right\rceil$. In other words, $C_{l}(x)$ is the union of all $k$ th level cylinder sets $\mathcal{C}_{k}$ such that $\operatorname{proj}(x) \in \operatorname{proj}\left(f_{\left(i_{1}, j_{1}\right)} \circ \cdots \circ f_{\left(i_{l}, j_{l}\right)}(\Lambda) \cap \mathcal{C}_{k}\right)$.


Like during the calculations of box dimension, $f_{\left(i_{l}, j_{l}\right)}^{-1} \circ \cdots \circ f_{\left(i_{1}, j_{1}\right)}^{-1}\left(C_{l}(x)\right)$ is the $k-l$ th level cylinder set of the IFS $\left\{g_{i}: x \mapsto \frac{x+i}{m}\right\}_{i \in Q}$, which contains $\operatorname{proj}\left(f_{\left(i_{l}, j_{l}\right)}^{-1} \circ \cdots \circ f_{\left(i_{1}, j_{1}\right)}^{-1}(x)\right)=\sum_{r=1}^{\infty} \frac{i_{l+r}}{m^{r}}$. That is,

$$
f_{\left(i_{1}, j_{1}\right)} \circ \cdots \circ f_{\left(i_{l}, j_{l}\right)}\left(\operatorname{proj}^{-1}\left[\sum_{r=1}^{k-l} \frac{i_{l+r}}{m^{r}}, \sum_{r=1}^{k-l} \frac{i_{l+r}}{m^{r}}+\frac{1}{m^{k-l}}\right]\right)=C_{l}(x)
$$

Therefore, by using the definition of $\mu$

$$
\mu\left(B_{\sqrt{2} / n^{l}}(x)\right) \geq \mu\left(C_{l}(x)\right)=\frac{t_{i_{1}}^{\frac{\log m}{\log n}-1} \cdots t_{i_{l}}^{\frac{\log m}{\log n}-1} \cdot t_{i_{l+1}}^{\frac{\log m}{\log n}} \cdots t_{i_{k}}^{\frac{\log m}{\log n}}}{\left(\sum_{i^{\prime} \in Q} t_{i^{\prime}}^{\frac{\log m}{\log n}}\right)^{k}}
$$

Thus,
$\frac{\log \mu\left(B_{\sqrt{2} / n^{l}}(x)\right)}{-l \log n} \leq \frac{\left\lceil l \frac{\log n}{\log m}\right\rceil \log m}{l \log n} \cdot \frac{\log \sum_{i^{\prime} \in Q} t_{i^{\prime}}^{\frac{\log m}{\log n}}}{\log m}+\frac{-1}{\log n}\left(\frac{1}{\left\lceil l \frac{\log n}{\log m}\right\rceil} \sum_{r=1}^{\lceil l \log n} \log m \quad \log t_{i_{r}}-\frac{1}{l} \sum_{r=1}^{l} \log t_{i_{r}}\right)$.
Hence, if

$$
\begin{equation*}
\left.\limsup _{l \rightarrow \infty} \frac{\left(\prod_{r=1}^{\lceil l \log n} \log \right\rceil}{} t_{i_{r}}\right)^{1 /\left\lceil l \frac{\log n}{\log m}\right\rceil}\left(\prod_{r=1}^{l} t_{i_{r}}\right)^{1 / l} \geq 1 \tag{8}
\end{equation*}
$$

then

$$
\liminf _{l \rightarrow \infty} \frac{\log \mu\left(B_{\sqrt{2} / n^{l}}(x)\right)}{-l \log n} \leq \frac{\log \sum_{i \in Q} t_{i}^{\frac{\log m}{\log n}}}{\log m} \text { for every } x \in \Lambda
$$

and by Lemma 1

$$
\begin{equation*}
\operatorname{dim}_{H} \Lambda \leq \frac{\log \sum_{i \in Q} t_{i}^{\frac{\log m}{\log n}}}{\log m} \tag{9}
\end{equation*}
$$

To show (8) holds, we need the following simple lemma:
Lemma 2. Let $\left\{a_{n}\right\}$ be a sequence of positive real numbers and let $c>1$. If $\lim \sup _{n \rightarrow \infty} \frac{a_{\lceil c n\rceil}}{a_{n}}<1$ then $\liminf _{n \rightarrow \infty} a_{n}=0$.

But for every sequence $i_{1}, i_{2}, \ldots, i_{r}, \ldots$ and every $l \geq 1$

$$
\left(\prod_{r=1}^{l} t_{i_{r}}\right)^{1 / l} \geq 1
$$

thus (9) holds.
Remark 3. Since we have shown that the self-affine measure $\mu$ with probabilities defined in (6) has maximal dimension, i.e.

$$
\operatorname{dim}_{H} \mu=\operatorname{dim}_{H} \Lambda
$$

by Remark 2 we get

$$
\operatorname{dim}_{B} \Lambda=\operatorname{dim}_{H} \Lambda \Leftrightarrow t_{i}=\frac{M}{N} \text { for every } i \in Q
$$

In particular, in that case the $\frac{\log \sum_{i \in Q} t_{i}^{\frac{\log m}{\log n}}}{\log m}$-dimensional Hausdorff measure is positive and finite.

