

# Digital Double Barrier Options: Several Barrier Periods and Structured Floors

**Sühan Altay**

( joint work with S. Gerhold, R. Haidinger and K. Hirhager )

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## Objectives:

- Determine the price of digital double barrier options with an arbitrary number of barrier periods.
- Find the exact value of the structured floor for structured notes whose individual coupons are digital double barrier options.
- Approximate the value of the structured floor by the price of a corridor option.

# Outline

- 1 Preliminaries and Pricing for One Period
- 2 Double Barrier Digitals with Finitely Many Periods
- 3 Structured Floors
- 4 Approximation by a Corridor Option

# Setup and Pricing for One Period

- The underlying  $(S_t)_{t \geq 0}$  has the risk-neutral dynamics

$$\frac{dS_t}{S_t} = r dt + \sigma dW_t,$$

with the initial value  $S_0 > 0$ , constant interest rate  $r \geq 0$ , volatility  $\sigma > 0$  and a standard Brownian motion  $W$ .

- At maturity  $T_0 + P$ , the payoff is one unit of currency if the underlying has stayed between the two barriers:

$$C_1 := \mathbb{I}_{\{B_{\text{low}} < S_t < B_{\text{up}}, \forall t \in [T_0, T_0 + P]\}}, \quad (1)$$

where  $T_0 > 0$ ,  $P > 0$  and  $B_{\text{up}} > B_{\text{low}} > 0$  are barriers.

# Setup and Pricing for One Period

- Let us denote the price of this “one-period double barrier digital” at  $t < T_0$  by

$$BD(S_t, t; \{T_0\}, P, B_{\text{low}}, B_{\text{up}}, r) := e^{-r(T_0+P-t)} \mathbf{E}[C_1 | \mathcal{F}_t]. \quad (2)$$

- The value function

$$f(S, t) := BD(S, t; \{T_0\}, P, B_{\text{low}}, B_{\text{up}}, r)$$

satisfies the Black–Scholes PDE

$$\frac{\partial f}{\partial t} + \frac{1}{2} \sigma^2 S^2 \frac{\partial^2 f}{\partial S^2} + rS \frac{\partial f}{\partial S} - rf = 0$$

with the terminal condition  $f(S, T_0 + P) = 1$ , for  $S \in (B_{\text{low}}, B_{\text{up}})$ , and the boundary conditions  $f(B_{\text{low}}, t) = f(B_{\text{up}}, t) = 0$  for  $t \in [T_0, T_0 + P]$ .

# Setup and Pricing for One Period

- Use the standard transformation  $f(S, t) = e^{\alpha x + \beta \tau} U(x, \tau)$ , where

$$\begin{aligned}x &:= \log(S/B_{\text{low}}), & \tau &:= \frac{1}{2}\sigma^2(T_0 + P - t), \\ \alpha &:= -\frac{1}{2} \left( \frac{2}{\sigma^2}r - 1 \right), & \beta &:= -\frac{2r}{\sigma^2} - \alpha^2,\end{aligned}\tag{3}$$

to transform the Black–Scholes PDE into the heat equation

$$\frac{\partial^2 U}{\partial x^2} = \frac{\partial U}{\partial \tau}.\tag{4}$$

- The boundary conditions in the new coordinates are

$$U(0, \tau) = U(L, \tau) = 0, \quad \tau \in [0, p],\tag{5}$$

where  $L := \log(B_{\text{up}}/B_{\text{low}})$  and  $p = \frac{1}{2}\sigma^2 P$ . The terminal condition translates to the initial condition

$$U(x, 0) = e^{-\alpha x}, \quad x \in (0, L).\tag{6}$$

## Proposition

For  $0 \leq t < T_0$ , the price of a barrier digital with barrier period  $[T_0, T_0 + P]$  and payoff  $C_1$  at  $T_0 + P$  is

$$BD(S_t, t; \{T_0\}, P, B_{\text{low}}, B_{\text{up}}, r) = \sqrt{2\pi} \left( \frac{S}{B_{\text{low}}} \right)^\alpha \sum_{k=1}^{\infty} k \frac{1 - (-1)^k e^{-\alpha L}}{\alpha^2 L^2 + k^2 \pi^2} e^{-(\frac{k\pi}{L})^2 \rho + \beta \tau} \\ \times \int_{-\frac{x}{\sqrt{2(\tau-\rho)}}}^{\frac{L-x}{\sqrt{2(\tau-\rho)}}} \sin \left( \frac{k\pi}{L} (x + y \sqrt{2(\tau-\rho)}) \right) e^{-y^2/2} dy. \quad (7)$$

- this is a *rear-end* barrier option, because the two barriers are alive only towards the end of the contract (see Hui (1997)) .
- in probabilistic terms, we are integrating the probability to stay between the barriers (see e.g Borodin and Salminen (2002)) , with  $S_{T_0}$  viewed as a parameter, against the law of  $S_{T_0}$ .

## Sketch of the Proof:

- First consider the rectangle  $(0, L) \times (0, p)$ , the (unique) solution can be found by separation of variables

$$U(x, \tau) = \sum_{k=1}^{\infty} b_k \sin\left(\frac{k\pi}{L}x\right) e^{-\left(\frac{k\pi}{L}\right)^2\tau}, \quad (x, \tau) \in (0, L) \times (0, p), \quad (8)$$

where

$$b_k := \frac{2}{L} \int_0^L e^{-\alpha x_1} \sin\left(\frac{k\pi}{L}x_1\right) dx_1 = 2k\pi \frac{1 - (-1)^k e^{-\alpha L}}{\alpha^2 L^2 + k^2 \pi^2}$$

are the Fourier coefficients of the boundary function  $U(x, 0) = e^{-\alpha x}$ .

- Next, at  $\tau = p$ , the solution is given by (8) for  $0 < x < L$  and vanishes otherwise.

$$U(x, p) = \begin{cases} \sum_{k=1}^{\infty} 2k\pi \frac{1 - (-1)^k e^{-\alpha L}}{\alpha^2 L^2 + k^2 \pi^2} \sin\left(\frac{k\pi}{L}x\right) e^{-\left(\frac{k\pi}{L}\right)^2 p}, & 0 < x < L, \\ 0, & x \leq 0 \text{ or } x \geq L. \end{cases} \quad (9)$$

- Finally solve for  $U$  in the region  $\mathbb{R} \times (p, \frac{1}{2}\sigma^2(T_0 + P))$ . A solution is found by convolving the initial condition (9) with the heat kernel (see Evans [4, p. 47]).



# Double Barrier Digitals with Finitely Many Periods

- $n$  tenor dates

$$0 < T_0 < \dots < T_{n-1}$$

and fixed period length  $P > 0$  satisfying  $T_i + P \leq T_{i+1}$  for  $i = 0, \dots, n-2$ .

- Consider a contract that pays one unit of currency at time  $T_{n-1} + P$ , if the underlying has remained between the two barriers  $B_{\text{low}}$  and  $B_{\text{up}}$  during each of the time intervals  $[T_i, T_i + P]$ ,  $i = 0, \dots, n-1$ .
- The price of this “multi-period double barrier digital” unless it is not knocked out before  $t$  is given by

$$BD(S_t, t; \{T_0, \dots, T_{n-1}\}, P, B_{\text{low}}, B_{\text{up}}, r) := e^{-r(T_{n-1}+P-t)} \mathbf{E} \left[ \prod_{i=1}^n C_i \middle| \mathcal{F}_t \right], \quad (10)$$

where

$$C_i := \mathbb{I}_{\{B_{\text{low}} < S_t < B_{\text{up}}, \forall t \in [T_{i-1}, T_{i-1} + P]\}}.$$

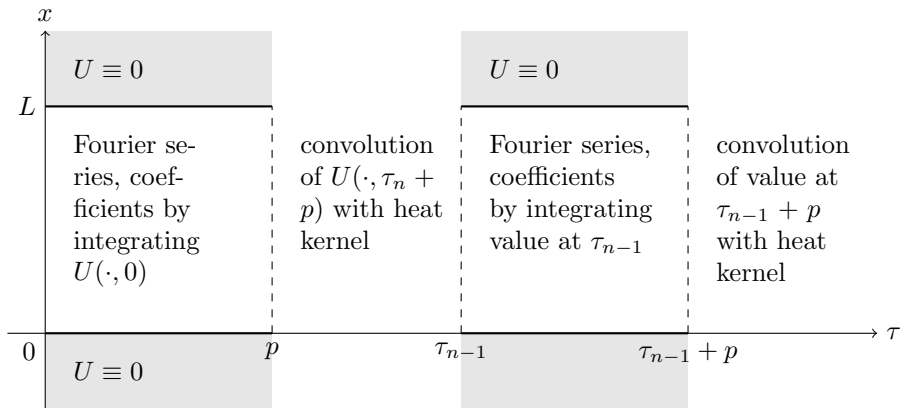


Figure: Solving the boundary value problem for an arbitrary number of barrier periods.

The  $n$  barrier periods  $[T_i, T_i + P]$  are mapped to  $[\tau_i, \tau_i + p]$ , where

$$\tau_i := \frac{1}{2}\sigma^2(T_{n-1} - T_{i-1}), \quad i = n, \dots, 1,$$

are the images of the barrier period endpoints under the coordinate change.

## Theorem

The value function (10) equals  $e^{\alpha x + \beta \tau} U(x, \tau)$ , where for  $j \in \{0, \dots, n-1\}$ ,  $\tau_{n-j} \leq \tau \leq \tau_{n-j} + p$ ,  $0 < x < L$ , we have

$$U(x, \tau) = \underbrace{\int_{-\infty}^{\infty} \dots \int_{-\infty}^{\infty}}_j \underbrace{\int_0^L \dots \int_0^L}_{j+1} \sum_{k_1=0}^{\infty} \dots \sum_{k_{j+1}=0}^{\infty} g_j(k_1, \dots, k_{j+1}; x_1, \dots, x_{j+1}; y_1, \dots, y_j; x, \tau) dx_1 \dots dx_{j+1} dy_1 \dots dy_j, \quad (11)$$

whereas for  $j \in \{0, \dots, n-1\}$ ,  $\tau_{n-j} + p < \tau < \tau_{n-(j+1)}$  (with  $\tau_0 := \infty$ ),  $x \in \mathbb{R}$ , we have

$$U(x, \tau) = \underbrace{\int_{-\infty}^{\infty} \dots \int_{-\infty}^{\infty}}_{j+1} \underbrace{\int_0^L \dots \int_0^L}_{j+1} \sum_{k_1=0}^{\infty} \dots \sum_{k_{j+1}=0}^{\infty} h_j(k_1, \dots, k_{j+1}; x_1, \dots, x_{j+1}; y_1, \dots, y_{j+1}; x, \tau) dx_1 \dots dx_{j+1} dy_1 \dots dy_{j+1}. \quad (12)$$

where  $h_j$  and  $g_j$  are defined recursively by

## Theorem

$$h_j(k_1, \dots, k_{j+1}; x_1, \dots, x_{j+1}; y_1, \dots, y_{j+1}; X, \tau)$$

$$:= \frac{1}{\sqrt{2\pi}} e^{-y_{j+1}^2/2} \mathbb{I} \left[ -\frac{x}{\sqrt{2(\tau - (\tau_{n-j} + p))}}, \frac{L-x}{\sqrt{2(\tau - (\tau_{n-j} + p))}} \right] (y_{j+1})$$

$$\cdot g_j(k_1, \dots, k_{j+1}; x_1, \dots, x_{j+1}; y_1, \dots, y_j; X + y_{j+1} \sqrt{2(\tau - (\tau_{n-j} + p))}, \tau_{n-j} + p)$$

and

$$g_j(k_1, \dots, k_{j+1}; x_1, \dots, x_{j+1}; y_1, \dots, y_j; X, \tau)$$

$$:= \frac{2}{L} \sin \frac{k_{j+1} \pi x_{j+1}}{L} \sin \frac{k_{j+1} \pi X}{L} e^{-(k_{j+1} \pi / L)^2 (\tau - \tau_{n-j})}$$

$$\cdot h_{j-1}(k_1, \dots, k_j; x_1, \dots, x_j; y_1, \dots, y_j; X_{j+1}, \tau_{n-j}),$$

with the recursion starting at

$$g_0(k_1; x_1; ; X, \tau) := \frac{2}{L} e^{-\alpha x_1} \sin \frac{k_1 \pi x_1}{L} \sin \frac{k_1 \pi X}{L} e^{-(k_1 \pi / L)^2 \tau}. \quad (13)$$

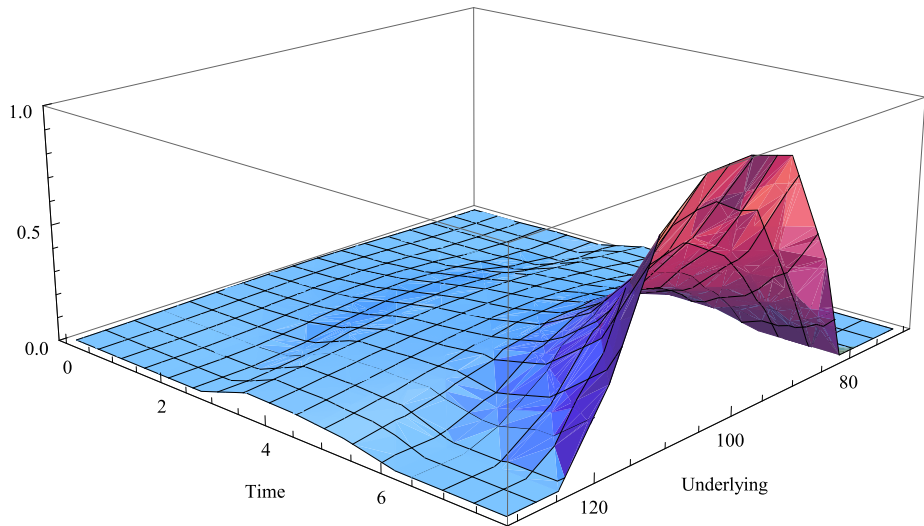
## Idea of the Proof:

- To iterate the argument of Proposition (see Figure 1)
- Using separation of variables in the barrier periods, and convolution with the heat kernel for the periods in between
- The required initial condition at the left boundary comes from the previous step of the iteration (for  $j = 0$  also from the payoff).

## Remarks:

- Proposition corresponds to (12) for  $j = 0$
- If a different option (a call, say) with the same barrier conditions is to be priced instead of a digital payoff, the quantity  $e^{-\alpha X_1}$  in (13) should be replaced by the appropriate payoff  $U(x_1, 0)$ .

# Numerical Implementation



**Figure:** Value function of a double barrier digital with two barrier periods with parameters  $r = 0.01$ ,  $\sigma = 0.15$ ,  $B_{\text{low}} = 80$ ,  $B_{\text{up}} = 120$ ,  $\{T_0, T_1\} = \{1, 6\}$ , and  $P = 2$ .

# Structured Notes and Floors

- Tenor structure satisfies  $T_{i-1} + P = T_i$  for  $i \in \{1, \dots, n-1\}$ , and define  $T_n := T_{n-1} + P$ .
- Consider a structured note with  $n$  coupons, where the  $i$ -th coupon consists of a payment of

$$C_i = \mathbb{I}_{\{B_{\text{low}} < S_t < B_{\text{up}}, \forall t \in [T_{i-1}, T_i]\}}, \quad i \in \{1, \dots, n\}, \quad (14)$$

at time  $T_i$ .

- In addition, the holder receives the terminal compensation

$$\left( F - \sum_{i=1}^n C_i \right)^+ \quad (15)$$

at  $T_n$ , where  $F > 0$ .

- These coupons can be priced by the Proposition (replace  $T_0$  by  $T_{i-1}$ ).

(Question): HOWEVER, HOW TO GET A HANDLE ON THE LAW OF  $A := \sum_{i=1}^n C_i$ ?

# Structured Floor

(Answer): THE LAW OF  $A := \sum_{i=1}^n C_i$  IS LINKED TO BARRIER OPTIONS WITH SEVERAL BARRIER PERIODS.

**Observation:** The moments

$$\mathbf{E}[A^\nu] = \sum_{i=0}^n i^\nu \mathbf{P}[A = i], \quad \nu \in \{1, \dots, n-1\}, \quad (16)$$

of  $A$  are linear combinations of multi-period double barrier option prices, with coefficients

$$c(\nu, J) := \sum_{\substack{0 \leq i_1, \dots, i_n \leq \nu \\ \text{supp}(\mathbf{i})=J}} \binom{\nu}{i_1, \dots, i_n}, \quad J \subseteq \{1, \dots, n\}. \quad (17)$$

(The notation  $\text{supp}(\mathbf{i}) = J$  means that  $J$  is the set of indices such that the corresponding components of the vector  $\mathbf{i} = (i_1, \dots, i_n)$  are non-zero.)



# Structured Floor

## Theorem

The price of the structured floor (15) at time  $t = 0$  can be expressed as

$$e^{-rT_n} \mathbf{E}[(F - A)^+] = e^{-rT_n} \sum_{i=0}^{n \wedge [F]} (F - i) \mathbf{P}[A = i], \quad (18)$$

where

$$\mathbf{P}[A = n] = BD(S_0, 0; \{T_0\}, T_n - T_0, B_{\text{low}}, B_{\text{up}}, 0). \quad (19)$$

The other point masses  $\mathbf{P}[A = i]$  in (18) can be recovered from the moments of  $A$  by solving (16) (including  $\nu = 0$ , of course). The moments in turn can be computed from barrier digital prices by ( $\nu \in \{1, \dots, n - 1\}$ )

$$\mathbf{E}[A^\nu] = \sum_{J \subseteq \{1, \dots, n\}} c(\nu, J) \cdot BD(S_0, 0; \{T_j : j \in J\}, P, B_{\text{low}}, B_{\text{up}}, 0), \quad (20)$$

where the coefficients  $c(\nu, J)$  are defined in (17).

## Proof.

The expression (18) is clear. The event in (19) means that all of the  $n$  coupons (14) are paid. By our assumption that  $T_i = T_{i-1} + P$ , its risk-neutral probability is the (undiscounted) price of a double barrier digital with one barrier period  $[T_0, T_n]$ , which yields (19). To prove (20), we calculate

$$\begin{aligned}\mathbf{E}[A^\nu] &= \mathbf{E}\left[\left(\sum_{i=1}^n C_i\right)^\nu\right] = \sum_{i_1, \dots, i_n} \binom{\nu}{i_1, \dots, i_n} \mathbf{E}[C_1^{i_1} \dots C_n^{i_n}] \\ &= \sum_{i_1, \dots, i_n} \binom{\nu}{i_1, \dots, i_n} \mathbf{E}\left[\prod_{\substack{j=1 \\ i_j > 0}}^n C_j\right] \\ &= \sum_{J \subseteq \{1, \dots, n\}} \left( \sum_{\substack{i_1, \dots, i_n \\ \text{supp}(\mathbf{i})=J}} \binom{\nu}{i_1, \dots, i_n} \right) \mathbf{E}\left[\prod_{j \in J} C_j\right].\end{aligned}$$

Now observe that  $\prod_{j \in J} C_j$  is the payoff of a double barrier digital with barrier periods  $[T_j, T_j + P]$  for  $j \in J$ . □

# Approximation of Structured Floor

- Numerical quadrature may be too involved for a large number of coupons.
- Let us fix a maturity  $T = T_n$  and assume that the  $n$  coupon periods

$$\mathcal{T}_i^n := ]\frac{i-1}{n}T, \frac{i}{n}T], \quad i \in \{1, \dots, n\},$$

have length  $T/n$ .

- For large  $n$ , the proportion of intervals during which the underlying stays inside the barrier interval

$$\mathcal{B} := [B_{\text{low}}, B_{\text{up}}]$$

is similar to the proportion of time that the underlying spends inside  $\mathcal{B}$ , i.e., the occupation time.

## Theorem

Let  $(S_t)_{t \geq 0}$  be a continuous stochastic process such that for each real  $c$  the level set  $\{t \geq 0 : S_t = c\}$  has a.s. Lebesgue measure zero. Then we have a.s.

$$\lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{I}_{\{S_t \in \mathcal{B} \ \forall t \in \mathcal{T}_i^n\}} = \frac{1}{T} \int_0^T \mathbf{1}_{\mathcal{B}}(S_t) dt.$$

Theorem suggests the approximation

$$e^{-rT} \mathbf{E}[(F - A)^+] \approx e^{-rT} \frac{n}{T} \mathbf{E} \left[ \left( \frac{FT}{n} - \int_0^T \mathbf{1}_{\mathcal{B}}(S_t) dt \right)^+ \right] \quad (21)$$

for the price of the structure floor (15). It is obtained from replacing  $F$  by  $F/n$  in the relation

$$\mathbf{E}[(nF - A)^+] \sim n \mathbf{E} \left[ \left( F - \frac{1}{T} \int_0^T \mathbf{1}_{\mathcal{B}}(S_t) dt \right)^+ \right], \quad n \rightarrow \infty.$$

which follows from above.

# Corridor Option

$$e^{-rT} \mathbf{E}[(F - A)^+] \approx e^{-rT} \frac{n}{T} \mathbf{E} \left[ \left( \frac{FT}{n} - \int_0^T \mathbf{1}_B(S_t) dt \right)^+ \right]$$

- On the right-hand side, we recognize the price of a put on the occupation time of  $(S_t)$ , also called a corridor option.
- Fusai (2000) studied such options in the Black–Scholes model. In particular, his Theorem 1 gives an expression for the Laplace transform of the characteristic function of  $\int_0^T \mathbf{1}_B(S_t) dt$ .
- By using numerical inversion techniques, Fusai (2000) shows how to price corridor options.

## Proof.

For  $1 \leq i \leq n$ , define processes  $(X_{ni}(t))_{0 \leq t \leq T}$  by

$$X_{ni}(t) := \begin{cases} 1 & \text{if } t \in \mathcal{T}_i^n \text{ and } S_u \in \mathcal{B} \forall u \in \mathcal{T}_i^n \\ 0 & \text{otherwise.} \end{cases}$$

Put  $X_n := \sum_{i=1}^n X_{ni}$ . We claim that, a.s., the function  $X_n(\cdot)$  converges pointwise on the set  $[0, T] \setminus \{t : S_t = B_{\text{low}} \text{ or } S_t = B_{\text{up}}\}$ , with limit  $\mathbf{1}_{\mathcal{B}}(S_t)$ . Indeed, if  $t \in [0, T]$  is such that  $S_t \notin \mathcal{B}$ , then  $X_n(t) = 0$  for all  $n$ . If, on the other hand,  $S_t \in \text{int}(\mathcal{B})$ , then  $t$  has a neighborhood  $V$  such that  $S_u \in \mathcal{B}$  for all  $u \in V$ , by continuity. Hence  $X_n(t) = 1$  for large  $n$ . Since we have pointwise convergence on a set of (a.s.) full measure, we can apply the dominated convergence theorem to conclude

$$\lim_{n \rightarrow \infty} \int_0^T X_n(t) dt = \int_0^T \mathbf{1}_{\mathcal{B}}(S_t) dt, \quad \text{a.s.}$$

But this is the desired result, since

$$\begin{aligned} \int_0^T X_n(t) dt &= \sum_{i=1}^n \int_0^T X_{ni}(t) dt \\ &= \sum_{i=1}^n \int_{\mathcal{T}_i^n} X_{ni}(t) dt = \sum_{i=1}^n |\mathcal{T}_i^n| \mathbf{1}_{\{S_t \in \mathcal{B} \forall t \in \mathcal{T}_i^n\}} = \frac{T}{n} \sum_{i=1}^n \mathbf{1}_{\{S_t \in \mathcal{B} \forall t \in \mathcal{T}_i^n\}}. \end{aligned}$$

# Numerical evaluation of the approximation

coupons	structure floor	corridor option	relative error
$n = 1$	7.63696	9.91563	0.2298060
$n = 2$	7.52979	9.24883	0.1858660
$n = 3$	7.42262	8.66698	0.1435750
$n = 4$	7.31545	8.06291	0.0927030
$n = 5$	7.20827	7.44886	0.0322987
$n = 6$	7.10110	7.31880	0.0297450
$n = 7$	6.99393	7.18558	0.0266714
$n = 8$	6.88677	7.03704	0.0213544
$n = 9$	6.77962	6.92288	0.0206936
$n = 10$	6.67232	6.80399	0.0193516

**Table:** Numerical evaluation of the approximation (21) with maturity  $T = 4$ , structure floor level  $F = 10$ , and  $n$  coupons. The other parameters are  $r = 0.01$ ,  $\sigma = 0.15$ ,  $B_{\text{low}} = 80$ , and  $B_{\text{up}} = 120$ .

- Approximation by a corridor option works only for period lengths tending to zero.
- One could also let the number of coupons tend to infinity for a fixed period length  $P$ , so that maturity increases linearly with  $n$ .

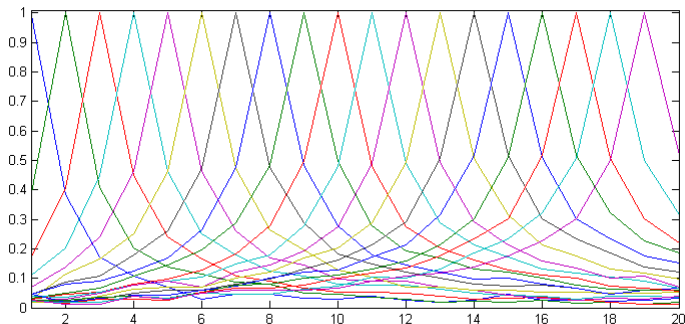


Figure: Correlation between individual coupons,  $n=20$



# Approximation of Sum of Coupons

- The correlation of the random variables  $C_i$  and  $C_j$  decreases for large  $|i - j|$ .
- Therefore, it is a natural question whether a "central limit theorem" holds, i.e., whether the sum of coupons  $A := \sum_{i=1}^n C_i$

$$\frac{A - \mathbf{E}[A]}{\sqrt{\mathbf{Var}[A]}}$$

converges in law to a standard normal random variable as  $n \rightarrow \infty$ .

- To have CLT (or certain normal approximation methods) we have to verify any of the "mixing" conditions (see Bradley (2005)).
- "Mixing" means, roughly, that random variables temporally far apart from one another are nearly independent.
- However, we could not verify for example, the  $\phi$  mixing condition.
- Numerical experiments also cast doubt on Gaussian limit law.

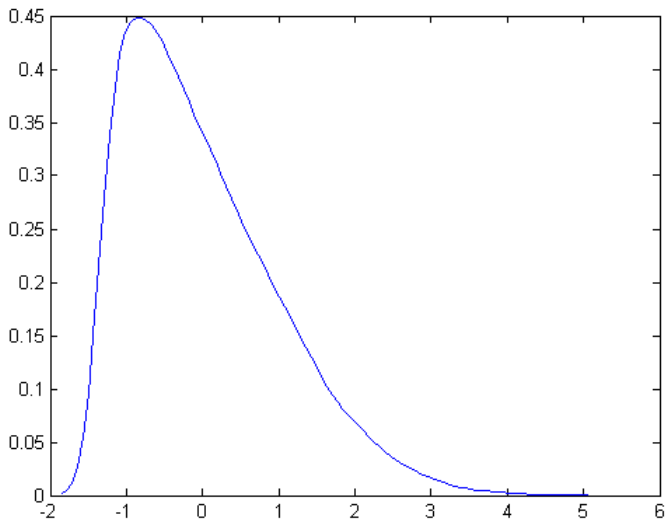








Figure: Density Approximation

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Thanks for listening...  
Comments! and Questions?

# Approximation of Sum of Coupons and $\phi$ -mixing

- On  $(\Omega, \mathcal{F}, \mathbf{P})$ , for any two  $\sigma$ -fields  $\mathcal{A}$  and  $\mathcal{B} \in \mathcal{F}$ ,
- $\phi(\mathcal{A}, \mathcal{B}) := \sup |\mathbf{P}(B|A) - \mathbf{P}(B)|$  given  $A \in \mathcal{A}$  and  $B \in \mathcal{B}$ , where the supremum is taken over all pairs of (finite) partitions  $\{A_1, \dots, A_i\}$  and  $\{B_1, \dots, B_j\}$  of  $\Omega$  such that  $A_i \in \mathcal{A}$  for each  $i$  and  $B_j \in \mathcal{B}$  for each  $j$ .
- Now suppose  $X := (X_k, k \in \mathbb{N})$  is a sequence of random variables. For  $0 \leq J \leq L \leq \infty$ , define the  $\sigma$ -field

$$\mathcal{F}_J^L := \sigma(X_k, J \leq k \leq L, (k \in \mathbb{N}))$$

and for each  $n \geq 1$  define

$$\phi(n) := \sup_{j \in \mathbb{N}} (\mathcal{F}_0^j, \mathcal{F}_{j+n}^\infty)$$

- The random sequence  $X$  is called  $\phi$ -mixing if  $\phi(n) \rightarrow 0$  as  $n \rightarrow \infty$ .
- In the context of our question, it can be shown that adapting the above notations in the obvious way,  $\phi(\sigma(W_s, s \leq t), \sigma(W_s, s \geq t + u)) = 1$ .